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Title: TURBO DECODER AND TURBO INTERLEAVER

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TURBO DECODER AND TURBO INTERLEAVER

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TECHNICAL FIELD

The present invention relates to a data processing system, and more particularly to a turbo decoder and a turbo interleaver.

BACKGROUND OF THE INVENTION

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Data signals, in particular those transmitted over a typically hostile RF interface (communication channel), are susceptible to error (channel noise) caused by interface. Various methods of error correction coding have been developed in order to minimize the adverse effects that a hostile interface has on the integrity of communicated data. This is also referred to as lowering the Bit Error Rate (BER), which is generally defined as the ratio of incorrectly received information bits to the total number of received information bits. Error correction coding generally involves representing digital data in ways designed to be robust with respect to bit errors. Error correction coding enables a communication system to recover original data from a signal that has been corrupted.

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Error correction code includes a convolution code, a parallel concatenated convolution code (so called turbo code). A convolution code transforms input sequence of bits into an output sequence of bits through the use of finite-state-machine, where additional bits are added to the data stream to allow for error-correction capability. In order to increase error-correction capability, the amount of additional bits added and the amount of memory preset in the finite-state-machine must be increased which increases decoding complexity.

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In the turbo coding system, a block of data may be encoded with a particular coding method resulting in systematic bits and two sets of parity bits. Additionally, the original block of input data may be rearranged with an interleaver and then encoded with the same method as that applied to the original input data. Encoded data (systematic bits and parity bits) are combined in some manner to form a serial bit stream and transmitted through the communication channel to a turbo decoding system. Turbo decoding systems operates on noisy versions of the systematic bits and the two sets of parity bits in two decoding stages to produce an estimate of the original

message bits. The turbo decoding system uses iterative decoding algorithms and consist of interleaver and deinterleaver stages, individually matched to constituent decoding stages. The decoding stages of the turbo decoding system uses BCJR algorithm, which was originally invented by Bahl, Cocke, Jelinek, and Raviv (hence
5 the name) to solve *maximum a posteriori probability* (MAP) detection problem. BCJR algorithm is a MAP decoder in that it minimizes the bit errors by estimating the a posteriori probabilities of the individual bits in a code word; to reconstruct the original data sequence, the soft outputs of the BCJR algorithm are hard-limited. The decoding stages exchange with each other the obtained soft output information and iteration of
10 decoding is ceased when a satisfactory estimate of transmitted information sequence has been achieved.

As the turbo code has extremely impressive performance which is very close to Shannon capacity limits, the 3G mobile radio systems such as W-CDMA and cdma2000 adopted them for channel coding.

15 3G wireless systems support the variable bit rate, which may result in full reconstruction of the turbo interleaver at every 10ms or 20ms frame. Accordingly, generating the whole interleaved address pattern at once consumes much time and requires a large-sized RAM to store the pattern.

Accordingly, high speed turbo interleaver which can support variable bit rate
20 and does not affect the performance of the turbo coder is required.

As is well-known, W-CDMA and cdma2000 are different in coding rate and interleaver. For example, coding rate of W-CDMA is 1/2, 1/3, 1/4 or 1/5 but coding rate of cdma2000 is 1/3, and frame size of the W-CDMA is one of twelve numbers 378, 370, 762, ..., and 20730 but that of the cdma2000 is arbitrary integer between 40
25 and 5114, and row of the block interleaver in W-CDMA is 32 (8-14 of them are unused) but that of the cdma2000 is 5, 10, or 20.

Accordingly, flexible and programmable decoders are required for 3G communication because global roaming is recommended between different 3G standards and the frame size may change on a frame base.

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SUMMARY OF THE INVENTION

Embodiments of the present invention provide an interleaver. The interleaver comprises a preprocessing means for preparing seed variables and an address

generation means for generating interleaved address using the seed variables on the fly. The seed variables are in forms of column vectors whose number of elements is equal to that of the rows of the two dimensional block interleaver. Consequently the number of seed variables is less than the size of block data. If the generated
5 interleaved address is larger than the size of block data, the generated interleaved address is discarded.

In some embodiments, the seed variables include a base column vector, an increment column vector and a cumulative column vector. The number of elements of all three column vectors is equal to the number of row of the interleaver block. The
10 cumulative column vector is updated by adding the increment vector to old cumulative column vector after interleaved addresses for one column is generated by adding the base column vector and the cumulative column vector. When updating the cumulative column vector, if elements of the updated cumulative column vector are larger than the number of column of the data block, the elements of the updated
15 cumulative column vector is subtracted by the number of the column of the data block.

Elements of the base column vector and the increment column vector are inter-row permuted.

Embodiments of the present invention provide a turbo decoding system. The turbo decoding system comprises an interleaver comprising a preprocessing means for
20 preparing seed variables and an address generation means for generating interleaved address using the seed variables, an address queue for storing the generated interleaved address equal to or smaller than the interleaver size, an SISO decoder performing recursive decoding and calculating log likelihood ratio, and an LLR memory connected to the SISO decoder and storing the log likelihood ratio, wherein
25 the SISO decoder accesses the input data and the log likelihood ratio alternately in a sequential order and in an interleaved order using the generated interleaved address.

In some embodiments, the generated interleaved address is reused as a write address for writing the log likelihood ratio outputted from the SISO decoder into the LLR memory.

30 Embodiments of the present invention provide a turbo decoding system comprising a processor for generating interleaved addresses and controlling hardware blocks, an address queue for storing the generated interleaved addresses, a buffer memory block including an LLR memory for storing log likelihood ratio and a

plurality of memory blocks for storing soft inputs, an SISO decoder connected to the buffer memory block, the SISO decoder including ACSA network for calculating the log likelihood ratio recursively from soft inputs and the log likelihood provided by the LLR memory and a plurality of memory blocks for storing intermediate results of the ACSA network.

In some embodiments, the processor preparing seed variables when the interleaver structure changes due to the change of the coding standard or bit rate, and generates the interleaved addresses column by column using the seed variables by simple add and subtract operations when the interleaved addresses are required.

In some embodiments, the SISO decoder supports Viterbi decoding mode. In Viterbi decoding mode, the ACSA network performs Viterbi recursion, the LLR memory stores traceback information outputted by the ACSA network, the processor processes traceback from the traceback information read from the LLR memory, and one of the memory of the SISO decoder stores path metric outputted by the ACSA network.

In some embodiments, the processor is a single-instruction and multiple-data (SIMD) processor. Preferably, the SIMD processor includes five processing elements, and wherein one of five processing elements controls the other four processing elements, processes scalar operation, and fetches, decodes, and executes instructions including control and multi-cycle scalar instructions, and wherein the other four processing elements only execute SIMD instruction.

Embodiments of the present invention provide an interleaving method for rearranging data block in a data communication system. The interleaving method comprises preparing seed variables and generating interleaved addresses on the fly using the seed variables.

BRIEF DESCRIPTION OF THE DRAWINGS

Other features of the present invention will be more readily understood from the following detailed description of the invention when read in conjunction with the accompanying drawings, in which:

Fig. 1 illustrates a basic turbo-encoding system.

Fig. 2 illustrates a typical eight-state RSC encoder of Fig. 1.

Fig. 3 shows block diagram of the turbo decoding system.

Fig. 4 shows extrinsic form of the turbo decoding system of Fig. 3.

Fig. 5 schematically shows block diagram of time-multiplex turbo decoding system according to an embodiment of the present invention.

5 Figs. 6A to 6C illustrate a simple example of prunable block interleaver with interleaver size $N = 18$ according to conventional interleaving technique.

Figs. 7A and 7B illustrate a simple example of prunable block interleaver with interleaver size $N=18$ according to the present invention.

Fig. 8 illustrate more specified block diagram of turbo decoding system of Fig. 5 in turbo decoding mode.

10 Fig. 9 illustrate more specified block diagram of turbo decoding system of Fig. 5 in Viterbi decoding mode.

Fig. 10 schematically shows detailed ACSA network and related memory blocks of Fig. 8.

15 Fig. 11 shows an ACSA unit contained in the ACSA A section 1022 for calculating forward metric $A_k(s)$ of Fig. 10.

Fig. 12 illustrates detailed SIMD processor of Fig. 8.

DETAILED DESCRIPTION OF THE INVENTION

20 The present invention now will be described more fully hereinafter with reference to the accompanying drawings, in which typical embodiments of the invention are shown. This invention may, however, be embodied in many different forms and should not be construed as limited to the embodiments set forth herein. Rather, these embodiments are provided so that this disclosure will be thorough and
25 complete, and will fully convey the scope of the invention to those skilled in the art.

Before proceeding to describe the embodiments of the present invention, typical turbo coding system will be described with reference to Figs. 1 to 4 for better understanding of the present invention.

30 Fig. 1 illustrates a basic turbo-encoding system 100 and Fig. 2 illustrates a typical eight-state RSC encoder 102, 106 of Fig. 1.

The encoder of the turbo coding system consists of two constituent systematic encoders 102, 106 joined together by means of an interleaver 104. Input data stream u is applied directly to first encoder 102, and the interleaved version of the input data

stream \mathbf{u} is applied to second encoder 106. The systematic bits (i.e., the original message bits) \mathbf{x}^s and the two sets of parity bits \mathbf{x}^{p1} and \mathbf{x}^{p2} generated by the two encoders 102, 106 constitute the output of the turbo encoding system 100 and are combined by a multiplexing means 108 to form a serial bit stream and transmitted
5 over the communication channel. Before transmission, puncturing may be performed if necessary.

The constituent encoder 102, 106 of the turbo encoding system is recursive systematic convolution (RSC) code, where one or more of the tap outputs in the shift-resistor $D1 \sim D3$ back to the input for obtaining better performance of the overall
10 turbo coding strategy.

Fig. 3 shows block diagram of the turbo decoding system 300. The turbo decoder 300 operates on noisy version of the systematic bits \mathbf{y}^s and noisy version of the two set of parity bits \mathbf{y}^{p1} and \mathbf{y}^{p2} . The turbo decoding system 300 uses iterative decoding algorithms and consist of interleaver 304 and deinterleaver 308, 310 stages,
15 individually matched to constituent decoding stages 302, 306. The systematic bits \mathbf{y}^s and first set of parity bits \mathbf{y}^{p1} of turbo encoded data are applied to the first SISO (Soft-Input-Soft-Output) decoder 302. Additionally, deinterleaved version of the metrics output of the second SISO decoder 306 is fed back to the first SISO decoder 302. The metrics output of the first SISO decoder 302 is applied to the second SISO decoder
20 306 via interleaver 304. The second set of parity bits \mathbf{y}^{p2} is applied to the second SISO decoder 306. The output of the deinterleaver 310 is applied to hard limiter 312 which outputs a bit stream of decoded data $\hat{\mathbf{u}}$ corresponding to the original raw data \mathbf{u} .

As stated earlier, the output of the deinterleaver 308 is fed back to the metrics input of the first SISO decoder 302. Thus the turbo decoding system 300 performs the
25 n th decoding iteration with an input metrics resulting from $(n-1)$ th decoding iteration. The total number of iteration is predetermined, or the iteration stops if a certain stopping criterion meets the qualification.

Fig. 4 shows extrinsic form of the turbo decoding system of Fig. 3, where I stands for interleaver, D for deinterleaver, and SISO for soft-input soft-output decoder,
30 which may use Log-MAP decoding algorithm, Max-Log-MAP, etc.

The first decoding stage produces a soft estimate $\Lambda_1(u_k)$ of systematic bit u_k expressed as log-likelihood ratio

$$\Lambda_1(u_k) = \log \left\{ \frac{P(u_k = 1 | \mathbf{y}^s, \mathbf{y}^{p1}, \Lambda_{2e}(\mathbf{u}))}{P(u_k = 0 | \mathbf{y}^s, \mathbf{y}^{p1}, \Lambda_{2e}(\mathbf{u}))} \right\}, k = 1, 2, \dots, N \quad (1)$$

where \mathbf{y}^s is the set of noisy systematic bits, \mathbf{y}^{p1} is the set of noisy parity bits generated by the first encoder 302, and $\Lambda_{2e}(\mathbf{u})$ is the extrinsic information about the set of message bits \mathbf{u} derived from the second decoding stage and fed back to the first stage.

Hence, the extrinsic information about the message bits derived from the first decoding stage is

$$\Lambda_{1e}(\mathbf{u}) = \Lambda_1(\mathbf{u}) - \Lambda_{2e}(\mathbf{u}) \quad \text{-----} (2)$$

where $\Lambda_{2e}(\mathbf{u})$ is to be defined.

Before application to the second decoding stage, the extrinsic information $\Lambda_{1e}(\mathbf{u})$ is reordered to compensate for the interleaving introduced in the turbo encoding system 100. In addition, the noisy parity bits \mathbf{y}^{p2} generated by the second encoder 106 are used as another input. Thus by using BCJR algorithm, the second decoding produces a more refined soft estimate of the message bits \mathbf{u} .

This estimate is de-interleaved to produce the total log-likelihood ratio $\Lambda_2(\mathbf{u})$. The extrinsic information $\Lambda_{2e}(\mathbf{u})$ fed back to the first decoding stage is therefore

$$\Lambda_{2e}(\mathbf{u}) = \Lambda_2(\mathbf{u}) - \Lambda_{1e}(\mathbf{u}) \quad \text{-----} (3)$$

where $\Lambda_{1e}(\mathbf{u})$ is itself defined by equation (2), and $\Lambda_2(\mathbf{u})$ is the log-likelihood ratio computed by the second decoding stage. Specifically, for the k th element of the vector \mathbf{u} , where we have

$$\Lambda_2(u_k) = \log \left\{ \frac{P(u_k = 1 | \mathbf{y}^s, \mathbf{y}^{p2}, \Lambda_{1e}(\mathbf{u}))}{P(u_k = 0 | \mathbf{y}^s, \mathbf{y}^{p2}, \Lambda_{1e}(\mathbf{u}))} \right\}, k = 1, 2, \dots, N \quad (4)$$

Through the application of $\Lambda_{2e}(\mathbf{u})$ to the first decoding stage, the feedback loop around the pair of decoding stages is thereby closed. Note that although in actual fact the set of noisy systematic bits \mathbf{y}^s is only applied to the first decoder 302 in Fig. 3, by formulating the information flow in the symmetric extrinsic manner depicted in Fig. 4 we find that \mathbf{y}^s is, in fact, also applied to the second decoding stage.

An estimate of message bits \mathbf{u} is computed by hard-limiting the log-likelihood ratio $\Lambda_2(\mathbf{u})$ at the output of the second decoding stage, as shown by

$\hat{\mathbf{u}} = \text{sgn}(\Lambda_2(\mathbf{u}))$, where the signum function operates on each element of $\Lambda_2(\mathbf{u})$ individually. To initiate turbo decoding algorithm, we simply set $\Lambda_{2e}(\mathbf{u}) = 0$ on the first iteration of the algorithm.

Now turbo decoding system of the present invention will be described.

5 Embodiments of the present invention provide multi-standard turbo decoding system with a processor on which software interleaver is run. Particularly, the present invention provides turbo decoding system with a configurable hardware SISO decoder and a programmable single-instruction and multiple-data (SIMD) processor performing flexible tasks such as interleaving. Software turbo interleaver is run on the
10 SIMD processor. The decoding system of the present invention can also support Viterbi decoding algorithm as well as turbo decoding algorithm such as Log-MAP algorithm and Max-Log-MAP algorithm.

The processor generates interleaved addresses for turbo decoder, supporting multiple 3G wireless standard at the speed of the hardware SISO decoder and
15 changing interleaver structure (i.e., frame size and bit rate) at a very short time with a small memory. To hide the timing overhead of interleaving changing, the interleaved addresses generation is split into two parts, pre-processing and incremental on-the-fly generation. Pre-processing part prepares a small number of seed variables and incremental on-the-fly generation part generates interleaved address based on the seed
20 variables on the fly. When bit rate changes, the processor carries out only the pre-processing part to prepare a small number of seed variables hence requiring short time and small memory. Whenever the interleaved address sequence is required, the processor generates the interleaved address using the seed variables. This splitting method reduces the timing overhead of interleaver and requires only a small memory
25 to save the seeding variables.

Fig. 5 schematically shows block diagram of time-multiplex turbo decoding system according to an embodiment of the present invention.

Turbo decoding system 500 of the present invention comprises an SISO decoder 502, processor 504 on which software interleaver is run and a Λ_e memory 506
30 for storing an extrinsic log-likelihood ratio (LLR) as illustrated in Fig. 5. Data are sequentially stored in Λ_e memory 506 as they are always read and written in-place. For each iteration, data are accessed in a sequential order for the *odd*th SISO decoding and

in an interleaved order for the *eventh* SISO decoding. Namely, in the *oddth* (first) SISO decoding, SISO decoder 502 receives data in a sequential order from the Λ_c memory 506 and calculate log-likelihood ratio, and the log-likelihood ratio is written into the Λ_c memory 506 in the sequential order. In the *eventh* (second) SISO decoding, SISO decoder 502 receives data in an interleaved order from the Λ_c memory 506 and calculate new log-likelihood ratio, and the new log-likelihood ratio is deinterleaved with the help of the address queue 508 and written into the Λ_c memory 506 in the sequential order. As shown with the dotted lines, the processor 504 provides interleaved address to read data in an interleaved order. The address queue 508 saves the addresses in an interleaved order so that the addresses can be used as the write addresses for Λ_c memory 506 when the SISO decoder 502 produces results after its latency. Accordingly, data in an interleaved order can be deinterleaved into the sequential order, and saved in the Λ_c memory 506 in a sequential order.

In addition to interleaving, the processor 504 can control the hardware blocks or interface with an external host, and processes the trellis termination and a stopping criterion during the first SISO decoding that does not need interleaved addresses. In Viterbi decoding mode, SISO decoder 502 is repeatedly used for the Viterbi recursion. The Λ_c memory 506 plays the roles of the traceback memory. The processor 504 processes the traceback from the traceback information read from the Λ_c memory 506.

Flexible software turbo interleaver run on the processor for turbo decoder will be described. To hide the timing overhead of interleaver change, the interleaved addresses generation is split into tow parts, pre-processing and incremental on-the-fly generation. Pre-processing part prepares a small number of seed variables and incremental on-the-fly generation part generates interleaved address based on the seed variables on the fly. When the interleaver size changes due to the change of bit rate or the communication standard itself, only the pre-processing part prepares a small number of seed variables, not all the interleaved address sequence. Through parallel processing using the seed variables, the processor generates interleaved addresses as fast as the hardware SISO decoding rate whenever the interleaved address sequence is required. The unit of on-the-fly address generation is a column of a block interleaver. Interleaved addresses are generated column by column, used for read addresses and stored in address queue, and then reused for write addresses for deinterleaving.

Before proceeding to describe interleaving technique according to the embodiment of the present invention, conventional interleaving techniques will be described for better understanding of the present invention.

The turbo decoders for wireless system are based on block turbo interleavers. Although the operations and parameters vary depending on the standards, W-CDMA and cdma2000 shares the prunable block interleaver structure where the interleaver is implemented by building a mother interleaver of a predefined size and then pruning unnecessary addresses. The mother interleavers can be viewed as two-dimensional matrices, where the entries are read in the matrix row by row and read out column by column. Before reading out the entries, intra- and inter-row permutations are performed.

Figs. 6A to 6C illustrate a simple example of a typical prunable block interleaver with interleaver size $N = 18$, where the data indexes are written in a matrix form. Incoming data are written into memory in two-dimensional matrix row by row as shown in Fig. 6A. Fig. 6B shows intra-row permutation data indexes from Fig. 6A. The intra-row permutation rule applied to this example is

$$y_{i,j} = b_i + [(j+1) * q_i] \bmod 5 \text{ ---- (5),}$$

where $y_{i,j}$ is the permuted index of the i th row and j th column, i and j are row and column indexes, respectively, $\mathbf{b} = (b_0, b_1, b_2, b_3) = (0, 5, 10, 15)$, and $\mathbf{q} = (q_0, q_1, q_2, q_3) = (1, 2, 3, 7)$. Fig. 6C shows inter-row permutation data result from Fig. 6B, which will be read out from the memory in column by column as a sequence of 17, 1, 13, 7, 2, 11, 9, ..., 0, 10, 5. The indexes 19, 18 exceeding the range of interest are pruned.

Now interleaving technique according to an embodiment of the present invention will be described with the example of Figs. 6A to 6C and with reference to Figs. 7A and 7B.

The present embodiment uses an increment vector \mathbf{w} of $w_i = q_i \bmod 5$ instead of \mathbf{q} , and a cumulative vector \mathbf{x}_j of

$$x_{i,j} = [(j+1) * q_i] \bmod 5 \text{ --- (6).}$$

Equation (6) can be rewritten as

$$y_{i,j} = b_i + x_{i,j} \text{ --- (7)}$$

and \mathbf{x}_j can be obtained recursively as

$$x_{i,j} = [(j+1) * (q_i \bmod 5)] \bmod 5$$

$$\begin{aligned}
&= [(j+1) * w_i] \bmod 5 \\
&= [(jw_i \bmod 5) + w_i] \bmod 5 \\
&= (x_{i,j-1} + w_i) \bmod 5, \text{ ----- (8)}
\end{aligned}$$

where $j = 1, 2, 3, 4, 5$ and $x_0 = w$.

5 As $0 \leq x_{i,j-1} < 5$ and $0 \leq w_i < 5$, $0 \leq x_{i,j-1} + w_i < 10$ and thus

$$\begin{aligned}
x_{i,j} &= x_{i,j-1} + w_i - 5 && \text{if } x_{i,j-1} + w_i \geq 5, \\
&= x_{i,j-1} + w_i && \text{otherwise} \quad \text{----- (9)}
\end{aligned}$$

According to the embodiments of the present invention, multiplication and modulo operations of equation (6) are replaced by cheaper operation, multiplication
10 by an addition and the modulo by a comparison and subtract operation.

As shown in Fig. 7A, \mathbf{b} , \mathbf{w} , and \mathbf{x}_0 for the first column of the block interleaver are calculated and stored in vector register of the processor in the preprocessing stage. The number of elements of each vector \mathbf{b} , \mathbf{w} , and \mathbf{x}_0 correspond to the number of elements of a column. Right side of the Fig. 7A shows that \mathbf{b} , \mathbf{w} , and \mathbf{x}_0 are stored in
15 the order of inter-row permutation such as $b_3, b_0, b_2, b_1 = (15, 0, 10, 5)$, and $\mathbf{w} = (w_3, w_0, w_2, w_1) = (2, 1, 3, 2) = \mathbf{x}_0$ in advance so as to free the on-the-fly generation from the inter-row permutation.

In the column by column on-the-fly address generation stage shown in Fig. 7B, the processor updates \mathbf{x}_j according to equation (9) and calculates the addresses based
20 on equation (7). Calculated addresses are sent to the address queue, if they are smaller than interleaver size N .

Referring to Fig. 7B, interleaved addresses for first column ($y_{i,1}$) is calculated by adding $\mathbf{b}_i + \mathbf{x}_0$. Since \mathbf{x}_0 is \mathbf{w} , $y_{i,1}$ is calculated by adding $\mathbf{b}_0 = (15, 0, 10, 5) + (2, 1, 3, 2) = (17, 1, 13, 7)$. After interleaved addresses for first column is calculated, \mathbf{x}_j is
25 updated by adding $\mathbf{x}_0 = (2, 1, 3, 2)$ and $\mathbf{w} = (2, 1, 3, 2)$. Thereby \mathbf{x}_1 is set to $(4, 2, 1, 4)$, where third element of \mathbf{x}_1 is 1 because $(3+3)$ is larger than 5 and thus subtracted by 5. Accordingly, interleaved addresses for second column ($y_{i,2}$) is calculated by adding $\mathbf{b}_i + \mathbf{x}_1 = (15, 0, 10, 5) + (4, 2, 1, 4) = (19, 2, 13, 7)$, where first element 19 is discarded since it is larger than or equal to the size $N (=18)$.

30 As described above, the present invention requires only small memory for storing the seed variables and performs add and subtract operations instead of modulo operation and multiplication operation.

The above example is a very simple one and the turbo interleavers in the real world such as W-CDMA and cdma2000 standard turbo interleavers are much more complex. However, the fundamental structure and the basic operations used in permutation rules are the same. The pseudocodes of the on-the-fly address generation for W-CDMA and cdma2000 are shown in Table 1 and Table 2, respectively.

In Table 1, C is the number of columns, R is the number of rows, p is a prime number, and $s(x)$ is a permutation sequence, which are determined from the interleaver size N according to the specification of the W-CDMA standard. Likewise C , R , and the binary power 2^n in Table 2 are determined from N according to the cdma2000 specification. The present invention handles those values also as seed variables and calculates them in advance at the preprocessing stage.

The on-the-fly generation flows of these real-world turbo interleavers are similar to the example. They also have base column vector \mathbf{b} , increment column vector \mathbf{w} , cumulative column vector \mathbf{x} , and a certain modulo base. \mathbf{x} is updated by adding \mathbf{w} to old \mathbf{x} value. If elements of the updated \mathbf{x} are larger than the modulo base, the elements of the updated cumulative column vector is subtracted by the modulo base. This operation substitutes a computationally expensive modulo operation. Then the interleaved addresses for one column are generated by adding \mathbf{b} and a vector that is calculated from \mathbf{x} .

Table 1

```

-----
0  column_counter = C-1
   loop:
25  1  x=x + w
     2  for each (i=0, 1,..., R-1) if( $x_i \geq p-1$ )  $x_i = x_i - (p-1)$ 
     3  load  $s(\mathbf{x})$  from the data memory
     4   $\mathbf{y} = \mathbf{b} + s(\mathbf{x})$ 
     5  for each (i=0, 1,..., R-1) if ( $y_i < N$ ) send  $y_i$  to the address queue
30  6  if((column_counter--) $\neq 0$ ) goto loop
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Table 2

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0  column_counter = C-1
5  loop:
    1  x=x + w
    2  for each (i=0, 1,..., R-1) if( $x_i \geq 2^n$ )  $x_i = x_i - 2^n$ 
    3  y = b + x
    4  for each (i=0, 1,..., R-1) if ( $y_i < N$ ) send  $y_i$  to the address queue
10  5  if((column_counter--) $\neq$  0) goto loop
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A SIMD (single-instruction and multiple-data) processor is suitable for this operation because the column is a vector and all the column entries get through exactly the same operations. However, in order to generate one address at every one or two cycles, some special instructions to make the long program short can be very helpful.

To speed up address generation, some customized instructions can be used to reduce the length of the loop in the on-the-fly generation part. The present invention introduces three processor instructions: STOLT (store to output port if less than), SUBGE (subtract if greater or equal), and LOOP. Each of them substitutes a sequence of three ordinary instructions but takes only one clock cycle to execute. For example, instruction STOLT corresponds to typical three RISC instructions SUB x, y, z; BRANCH if z \geq 0; STO x. Likewise, SIMD instruction SUBGE corresponds to three RISC instructions SUB x, y, z; BRANCH if z < 0; MOVE z, x .

Pruning can be mapped to STOLT. The function of STOLT is to send the calculated interleaved address to the address queue only if calculated interleaved address is smaller than N , which is needed for the pruning as in line 5 of the pseudocode of Table 1 and line 4 of Table 2.

Another conditional instruction SUBGE, is quite useful for the block interleavers that commonly use modulo operations. Instruction SUBGE substitutes a

modulo or remainder operation $a \bmod b$ if the condition $0 \leq a < 2b$ is satisfied, which corresponds to (9) and line 2 of Table 1 and Table 2.

Adopted in several DSP processors to reduce the loop overhead of the address generation, LOOP instruction is also helpful in our application. This instruction
5 conforms to a sequence of CMP, BNE (branch if not equal), and SUB instructions, which at once decrements the loop count and branches.

Using these special instructions, the present invention can reduce the lengths of the on-the-fly generation program loop of W-CDMA, cdma2000, and CCSDS to six, five, and four instructions, respectively. Using loop-unrolling technique, the
10 present invention can shorten further the loop length of the on-the-fly address generation parts by almost one instruction.

In the turbo interleaver pseudocodes of Table 1 and Table 2, each line corresponds to an instruction of the SIMD processor code. In Table 1, the line 2 corresponds to SUBGE, the line 5 to STOLT, and the line 6 to LOOP. The SUBGE
15 safely substitutes $x_i = x_i \bmod (p-1)$ because the condition $0 \leq x_i < 2(p-1)$ is satisfied ($0 \leq x_i < p-1$ and $0 \leq w_i < p-1$ before they are added). If $R = 10$ or 20 and the processor can process five data at once, lines 1–5 are repeated twice or four times to produce an entire column of the interleaver matrix. Similarly, in Table 2 the line 2 corresponds to SUBGE, the line 4 to STOLT, and the line 5 to LOOP.

20 Fig. 8 illustrates a block diagram of decoding system which can support turbo decoding and Viterbi decoding. Fig. 8 schematically shows data flow and address flow in turbo decoding mode. In Fig. 8, solid line indicates data flows and dotted line indicates address flow. Fig. 9 schematically shows data flow and address flow in Viterbi decoding mode.

25 Referring to Fig. 8, turbo decoding system of the present invention comprises SISO decoder 810, processor block 830, buffer memory block 850 and address queue 870. The processor block 830 includes SIMD (single-instruction and multiple-data) processor 832, an instruction memory 834, and a data memory 836. The SISO decoder 810 implements a typical sliding-window decoder. It includes ACSA (add-compare-selector-add) network 812, and plurality of memory blocks, Γ_1 memory 814, Γ_2
30 memory 816, Γ_3 memory 818, A memory 820 and hard decision memory 822. The buffer memory block 850 includes a Λ_e memory 852 for storing an extrinsic log-

likelihood ratio (LLR), plurality of memories for storing soft input data, y^s memory 854 for storing a nosy systematic bits multiplied by the channel reliability 862, $y^{p1} \sim y^{p3}$ memories 856, 858, 860 for storing a nosy parity bits multiplied by the channel reliability 862.

5 The ACSA network 812 calculates forward metrics A_k , backward metrics B_k , and extrinsic log-likelihood ratio (LLR). Memory blocks, Γ_1 memory 814, Γ_2 memory 816, Γ_3 memory 818 stores input data and memory block A 820 temporarily stores the calculated forward metrics. Hard decision output of the ACSA network 812 is stored in the hard decision memory 822. SIMD processor 832 also calculates stopping
10 criterion during SISO decoding from information stored in the hard decision memory 822. Input data are read into one of the memory blocks, Γ_1 memory 814, Γ_2 memory 816, Γ_3 memory 818 and used three times for calculating forward metric A_k , backward metric B_k , and LLR $\Lambda(u_k)$.

Software interleaver is run on the SIMD processor 832. As described earlier,
15 The SIMD processor 832 generates interleaved addresses column by column and the address queue 870 saves the interleaved addresses. When the SIMD processor 832 calculates interleaved read addresses, the address queue 870 whose length is the SISO latency saves the interleaved addresses in order to use them again as the write addresses into the Λ_e memory 852. Namely, when the ACSA network 812 produces
20 results using the data read from the interleaved addresses, the results are stored into the corresponding place of the Λ_e memory 852 with the write address stored in the address queue 870. In addition to the interleaving, the SIMD processor 832 controls the hardware blocks, interfaces with an external host, and processes the trellis termination and a stopping criterion during SISO decoding that does not need an
25 interleaver.

Since Viterbi algorithm does not calculate backward metrics in Viterbi decoding mode, some components of Fig. 8 are unused as shown in Fig. 9. In Fig. 9, components illustrated by dotted line such as address queue 970, channel reliability multiplication 962, and three Γ memories blocks 914, 916, 918 are not used in
30 Viterbi decoding mode. The Λ_e memory 852 of Fig. 8 serves as traceback memory 952 and the A memory 820 of Fig. 8 serves as path metric memory. The SIMD processor

932 processes the traceback from the traceback information read from traceback memory 952. SISO network 912 is used for the Viterbi forward trellis recursion.

Fig. 10 schematically shows detailed ACSA network and related memory blocks of Fig. 8. The SISO decoder 1000 in Fig. 10 implements sliding window turbo decoding technique in hardware. The ACSA network 1010 includes multiplexers 1012, 1014, 1016, 1018, 1020 and ACSA A 1022, two ACSA B 1024, 1026, and ACSA A 1028. ACSA A 1022 and ACSA B 1024 comprise eight ACSA units and ACSA A 1028 comprises fourteen ACSA units. An exemplary ACSA unit of ACSA A 1002 is illustrated in Fig. 11.

After input data sequences are stored in the Γ_1 memory 1050 and Γ_2 memory 1070, SISO decoding starts while new interleaved data sequences are stored in the Γ_3 memory 1090.

To implement the sliding window, first a window size of L input data are stored in Γ memories 1050, 1070, 1090. According to the MAP algorithm, this block of written values is read three times. These three operations are performed in parallel in each ACSA sections in Fig. 10. The ACSA B section 1026 implements sliding window algorithm, which require a dummy backward recursion of depth L . In order to avoid the use of a multiple-port memory, three separated Γ memories of depth L are used so that the defined operations operate on them in a cyclic way. A memory 1030 temporarily stores the calculated $A_k(s)$'s.

SISO outputs are obtained in the reversed order, but the correct order can be restored by properly changing the interleaving law of the decoder.

To support multiple standards, configurable ACSA units can be employed in the SISO decoder in Fig. 10. The ACSA unit shown in Fig. 11 can be adapted to various RSC codes of different constraint length K , different coding rate of 1/2 to 1/5, and/or arbitrary generator polynomial.

The Log-MAP algorithm outperforms the Max-Log-MAP algorithm if the channel noise is properly estimated. However, it is reported that the Max-Log-MAP is more tolerant to the channel estimation error than the Log-MAP algorithm. Thus present invention provides the ACSA unit that can be selected to use Log-MAP or Max-Log-MAP algorithm.

Fig. 11 shows an ACSA unit contained in the ACSA A section 1022 for calculating forward metric $A_k(s)$ of Fig. 10.

Forward metric $A_k(s)$ is calculated by equation (10)

$$A_k(s) = \ln \left(\sum_{s'} \exp [A_{k-1}(s') + \Gamma_k(s', s)] \right) \quad \text{----- (10)}$$

$$\approx \max_{s'} (A_{k-1}(s') + \Gamma_k(s', s))$$

5 where $A_k(s')$ is the forward metric of state s' at the previous time stamp $k-1$ and $\Gamma_k(s', s)$ is the logarithmic branch probability tat the trellis state changes from s' to s at time stamp k . The ACSA unit 1100 includes input multiplexers 1101, 1103, 1105, 1107, 1109, 1111, two adder blocks (CSA) 1113, 1115, two input adders 1117, 1119, comparator 112, lookup table 1123, output multiplexer 1127, and output adder
10 1127.

Two adder blocks (CSA) 1113, 1115 calculate branch probabilities $\Gamma_k(s', s)$ given by equation (11)

$$\Gamma_k(s', s) = \frac{1}{2} \left(x_k^s [\Lambda_e(u_k) + L_c y_k^s] + L_c \sum_{l=1}^n y_l' x_k^l \right) \quad \text{----- (11)}$$

where $\Lambda_e(u_k)$ is the extrinsic LLR information from the previous SISO
15 decoding and L_c is channel reliability. Input data $\Lambda(u_k) + L_c y_k^s$, $L_c y_k^{p0}$, and $L_c y_k^{p1}$ are selected by multiplexers 1101, 1103, 1105, 1107, 1109, 1111 which can change the coding rate and the transfer function. The input adder 1117 adds output branch probabilities $\Gamma_k(s_0', s)$ of the adder block 1113 and incoming data $A_{k-1}(s_0')$. The input adder 1119 adds output branch probabilities $\Gamma_k(s_1', s)$ of the adder block 1115
20 and incoming data $A_{k-1}(s_1')$. The comparator 1121 receives two input from the input adders 1117, 1119 and outputs max value of the two inputs and differential value between two inputs. Differential value is used to look up the table 1123 that stores approximation offsets of (10), and max value is transferred in the output adder 1127. The output multiplexer 1125 selects decoding algorithm, Log-MAP or Max-Log MAP.
25 If 0 is selected at the multiplexer 1125, output $A_k(s)$ of the adder 1127 is given as $A_k(s) = \max(A_{k-1}(s') + \Gamma_{k-1}(s', s))$, which corresponds to Max-Log MAP algorithm. Otherwise if the lookup table value is selected, output $A_k(s)$ of the adder 1127 is given

as $A_k(s) = \ln(\exp [A_{k-1}(s') + \Gamma_{k-1}(s', s)])$, which is compensated by the offset read from the lookup table and corresponds to Log-MAP algorithm.

In conventional hardware SISO decoders, the calculation of $\Gamma_k(s', s)$ is fixed, as all the x_k^l 's are fixed for the target turbo code. However, the present invention can
5 change the x_k^l values in equation by configuring the input multiplexers in Fig. 11. This allows the change of coding rate and generator polynomials of the RSC encoder. The ACSA unit in the figure can support the rate of 1/2 to 1/5 turbo codes with arbitrary generator polynomials. To support less coding rate, the input to $\Gamma_k(s', s)$ calculation logic should be increased. To support multiple constraint length K , the number of
10 ACSA units and interconnection between the units can be changed. As mentioned above, the multiplexer 1125 in the right determines the decoding algorithm: Log-MAP or Max-Log-MAP. If the reliable channel estimation from external host which calculates the channel estimation by using is obtained, for example, power control bit of the 3G communication systems, we can obtain better performance with Log-MAP
15 algorithm by setting the multiplexer passes the look-up table value. On the other hand, if nothing is known about the channel, the Max-Log-MAP is used to avoid error due to channel misestimation error by passing 0 to the final adder of the ACSA unit.

To keep pace with the hardware SISO described with reference to Figs. 10 and 11, it is preferable that interleaved address generation be parallel processing. The
20 present invention employs SIMD architecture because it is suitable for the simple and repetitive address generation and has simpler control and lower power consumption than superscalar or very long instruction word (VLIW) processor. Fig. 12 illustrates detailed SIMD processor of Fig. 8. In Fig. 12, dotted line indicates control flow and solid line indicates data flow. Considering the number of rows of W-CDMA block
25 interleaver is multiple of five, the SIMD processor 1200 includes five processing elements PE0 ~ PE4. The bit widths of instructions and data are 16. The first processing element PE0 (1201) controls the other four processing elements PE1 ~ PE4 and in addition it processes scalar operation. The first processing element PE0 1201 fetches, decodes, and executes instructions including control and multi-cycle scalar
30 instructions while the other four processing elements PE1 ~ PE4 only execute SIMD instructions. Instruction corresponding to program counter (PC) 1210 is fetched from instruction memory 1227 and temporarily stored into the instruction register (IR) 1213.

The fetched instruction is decoded by the decoder 1214 and then decoded instruction is executed accordingly in each the processing elements PE0 ~ PE4, for example, each ALU 1221 executes add operations. After execution of instruction is completed, program counter (PC) 1210 is incremented by PC controller 1211 and new instruction is fetched from the instruction memory 1227. The other four processing elements PE1 (1203), PE2 (1205), PE3 (1207), PE4 (1209) executes SIMD instructions. All the processing elements PE0 ~ PE4 includes resister block 1215 for storing data for parallel operations. The register block 1215 includes vector resisters VR0 ~ VR15 (1217). The register block 1215 of the first processing element PE0 also includes additional scalar resisters 1219 to store scalar and control data. The second and fifth processing elements PE1 ~ PE4 includes register 1225 for temporarily storing encoded instruction. SIMD instruction is not executed in all processing elements at the same time, but executed in one processing element after another so that a data memory port and I/O port can be shared in a time-multiplexed fashion, which saves memory access power and provides a simple I/O interface.

As mentioned before, specialized SIMD processor instructions, STOLT, SUBGE, and LOOP, are employed to replace common instruction sequences of three typical RISC instructions appearing in turbo interleaver programs. STOLT and SUBGE are SIMD instructions, whereas LOOP is a scalar control instruction, which is executed only in PE0.

According to at least one embodiment of the present invention, the two stages of the block turbo interleaver performs preprocessing and on-the-fly generation in the SIMD processor. The speed of on-the-fly generation dramatically improves using the SIMD processor because of its SIMD parallel processing capability and its support of three special instructions.

It should be noted that many variations and modifications may be made to the embodiments described above without substantially departing from the principles of the present invention. All such variations and modifications are intended to be included herein within the scope of the present invention, as set forth in the following claims.